DISCRETE DIFFERENTIAL GEOMETRY: AN APPLIED INTRODUCTION Keenan Crane • CMU 15-458/858



LECTURE 11: DISCRETE CURVES



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Curves, Surfaces, and Volumes

- In general, differential geometry studies *n*-dimensional manifolds; we'll focus mostly on low dimensions: curves (*n*=1), surfaces (*n*=2), and volumes (*n*=3)
- Why? Geometry we encounter in "every day life" (Common in applications!) • Low-dimensional manifolds are not baby stuff! :-)
- - *n*=1: unknot recognition (open as of July 2017)
 - *n*=2: Willmore conjecture (2012 for genus 1)
 - *n*=3: Geometrization conjecture (2003, \$1 million)
- Serious intuition gained by studying low-dimensional manifolds
- Conversely, problems involving very high-dimensional manifolds (e.g., statistics/ machine learning) involve less "deep" geometry than you might imagine!
 - fiber bundles, Lie groups, curvature flows, spinors, symplectic structure, ...
- Moreover... curves and surfaces are beautiful! (And sometimes boring for large *n*...)



Curves & Surfaces



*Or solids... but the boundary of a solid is a surface!

• Much of the geometry we encounter in life well-described by *curves* and *surfaces**



Smooth Descriptions of Curves & Surfaces

- Many ways to express the geometry of a curve or surface:
 - height function over tangent plane
 - local parameterization
 - Christoffel symbols coordinates / indices
 - differential forms "coordinate free"
 - moving frames change in *adapted frame*
 - Riemann surfaces (*local*); Quaternionic functions (*global*)
- We'll dive deep into one description (**differential forms**) and touch on others



• People can get very religious about these different "dialects"... best to be multilingual!

Discrete Descriptions of Curves & Surfaces

- Also *many* ways to discretize a surface
- For instance:
 - **implicit** *e.g.*, zero set of scalar function on a grid
 - good for changing topology, high accuracy
 - expensive to store / adaptivity is harder
 - hard to solve sophisticated equations *on* surface
 - explicit *e.g.*, polygonal surface mesh
 - changing topology, high-order continuity is harder
 - cheaper to store / adaptivity is much easier
 - more mature tools for equations *on* surfaces

• Don't be "religious"; use the right tool for the job!



explicit



implicit

Curves & Surfaces – Overview

- of view.
- <u>Smooth setting:</u>
 - express geometry via differential forms
 - will first need to think about *vector-valued* forms
- Discrete setting:
 - use explicit mesh as domain
 - express geometry via discrete differential forms
- **Payoff:** will become very easy to switch back & forth between smooth setting (scribbling in a notebook) and discrete setting (running algorithms on real data!)



• Goal: understand curves & surfaces from complementary smooth and discrete points





Discrete Curves

Discrete Curves in the Plane

i.e., a sequence of points connected by straight line segments:



• We'll define a **discrete curve** as a *piecewise linear* parameterized curve,



$$\begin{array}{c} & S \\ \hline & \\ \cdot & \\ \end{array} \\ s_n = L \end{array}$$

Shorthand: $\gamma_i := \gamma(s_i)$

Discrete Curves in the Plane—Example

• A simple example is a curve comprised of two segments:

$$\gamma(s) := \begin{cases} (s,0), & 0 \le s \le \\ (1,s-1), & 1 \le s \le \end{cases}$$



Discrete Curves and Discrete Differential Forms

- Equivalently, a discrete curve is determined by a discrete, *Rⁿ*-valued 0-form on a manifold simplicial 1-complex
- The 0-form values give the location of the vertices; interpolation by Whitney bases (hat functions) gives the map from each edge to R^n



$K = \{ (v_0, v_1), (v_1, v_2), (v_2, v_3), \}$ $(v_0), (v_1), (v_2), (v_3), \emptyset$

$\gamma(v_0)$	—	(33,66)
$\gamma(v_1)$	—	(79,36)
$\gamma(v_2)$	—	(118, 58)
$\gamma(v_3)$		(134, 47)



Differential of a Discrete Curve

- •We can now directly translate statements about **smooth** curves expressed via smooth exterior calculus into statements about discrete curves expressed using discrete exterior calculus
- •Simple example: the *differential* just becomes the edge vectors:







Discrete Tangent

tangents, yielding a vector per edge*



 $T(s) := d\gamma(\frac{d}{ds}) / |d\gamma(\frac{d}{ds})|$

*And no definition of the tangent at vertices!

• As in smooth setting, can simply normalize differential to obtain



 $T_{ij} := (d\gamma)_{ij} / |(d\gamma)_{ij}|$



Discrete Normal

planar curve as a 90-degree rotation of the (discrete) tangent:



• As in the smooth setting, we can express the (discrete) normals of a



 $N_{ij} = \mathcal{J}T_{ij}$

Regular Discrete Curve / Discrete Immersion

- •Recall that a smooth curve is *regular* if its differential is nonzero; this condition helps avoid "bad behavior" like sharp cusps
- •For a discrete curve, a nonzero differential merely prevents zero edge lengths; need something stronger to get "nice" curves
- •In particular, a *regular discrete curve* or *discrete immersion* is a discrete curve that is a **locally injective map**
- •Rules out zero edge lengths and zero angles





Discrete Curvature

•For a regular discrete curve, discrete curvature has several definitions

TURNING ANGLE



STEINER FORMULA

 $\gamma - \varepsilon N$





determined by its edge lengths and turning angles.

Q: Given only this data, how can we recover the curve?

 $\varphi_{i+1,i+2}$ **A:** Mimic the procedure from the smooth setting: Sum curvatures to get angles: $\varphi_{i,i+1} := \sum \theta_k$ Evaluate unit tangents: $T_{ij} := (\cos(\varphi_{ij}), \sin(\varphi_{ij}))$ Sum tangents to get curve: $\gamma_i := \sum \ell_{k,k+1} T_{k,k+1}$ k=1

Q: Rigid motions?

Fundamental Theorem of Discrete Plane Curves

- **Fact.** Up to rigid motions, a regular discrete plane curve is uniquely



Discrete Whitney Graustein

- If we adopt the definition of a discrete regular curve as one that is *locally injective*, then there is a discrete version of Whitney-Graustein that exactly mirrors the smooth one
- Has been carefully studied from several perspectives:
 - Constructive algorithm (case analysis) by Mehlhorn & Yap (1991)
 - Much simpler argument by Pinkall in terms of convex polyhedron: https://bit.ly/2BFtywA
- Both use powerful idea from (discrete) differential geometry: to find a "path" connecting two objects, find path from both objects to a canonical one, then compose... (uniformization, Delaunay, ...)

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<u>Leave a reply</u>

CONSTRUCTIVE WHITNEY-GRAUSTEIN TI The Discrete Whitney-Graustein Theorem OR HOW TO UNTANGLE CLOSED PLANAL

KURT MEHLHORN[†] AND CHEE-KENG YA

Abstract. The classification of polygons is considered in which two pol if one can be continuously transformed into the other such that for each adjacent edges overlap. A discrete analogue of the classic Whitney-Graustein that the winding number of polygons is a complete invariant for this classific constructive in that for any pair of equivalent polygons, it produces some sequ taking one polygon to the other. Although this sequence has a quadratic num be described and computed in real time.

Key words. polygons, computational algebraic topology, computational theorem, winding number

Let us consider regular closed discrete plane curves γ with n vertices and tangent winding number m. We assume that the length of γ is normalized to some arbitrary (but henceforth fixed) constant L. Up to orientation-preserving rigid motions such a γ is uniquely determined by a point

$$(\ell_1,\ldots,\ell_n,\kappa_1,\ldots,\kappa_n)\in (0,\infty)^n imes(-\pi,\pi)^n$$

satisfying

$$\ell_1+\ldots+\ell_n=L$$
 $\kappa_1+\ldots+\kappa_n=2\pi m$
 $\ell_1e^{ilpha_1}+\ldots+\ell_ne^{ilpha_n}=0$

where

$$\alpha_j = \kappa_1 + \ldots + \kappa_j.$$

fixed $(\kappa_1,\ldots,\kappa_n)\in imes(-\pi,\pi)^n$ satisfying Consider а $\kappa_1 + \ldots + \kappa_n = 2\pi m$ for some $m \in \mathbb{Z}$ and define $\alpha_1, \ldots, \alpha_n$ as above. Then the set of $(\ell_1,\ldots,\ell_n)\in (0,\infty)^n$ satisfying









Curvature Flow

Curvature Flow on Curves

- A curvature flow is a time evolution of a curve (or surface) driven by some function of its curvature.
- Such flows model physical *elastic rods*, can be used to find shortest curves (geodesics) on surfaces, or might be used to smooth noisy data (e.g., image contours).
- Two common examples: *length*shortening flow and elastic flow.





Discretizing a Gradient Flow

- Two possible paths for discretizing any gradient flow:
 - 1. **First** derive the gradient of the objective in the smooth setting, then discretize the resulting evolution equation.
 - 2. **First** discretize the objective itself, **then** take the gradient of the resulting discrete objective.
- •In general, *will not* lead to the same numerical scheme/algorithm!





(Does **NOT** commute in general.)

Length Shortening Flow

- The objective for length shortening flow is simply the total length of the curve; the flow is then the (L^2) gradient flow.
- For closed curves, several interesting features (Gage-Grayson-Hamilton):
 - •Center of mass is preserved
 - Curves flow to "round points"
 - Embedded curves remain embedded

 $\operatorname{length}(\gamma) := \int_{0}^{L} \left| \frac{d}{ds} \gamma \right| \, ds$ $\frac{d}{dt}\gamma = -\nabla_{\gamma} \text{length}(\gamma)$ 0.015

credit: Sigurd Angenent



Length Shortening Flow

Let length(γ) denote the total length of a regular plane curve $\gamma : [0, L] \to \mathbb{R}^2$, and consider a variation $\eta : [0, L] \to \mathbb{R}^2$ vanishing at endpoints. One can then show that

$$\frac{d}{d\varepsilon}|_{\varepsilon=0} \operatorname{length}(\gamma + \varepsilon\eta) =$$

Key idea: quickest way to reduce length is to move in the direction κN .

 $-\int_{0}^{L} \langle \eta(s), \kappa(s) N(s) \rangle \, ds$ $+ \epsilon \eta$

Length Shortening Flow—Forward Euler

- At each moment in time, move curve in normal direction with speed proportional to curvature
- "Smooths out" curve (e.g., noise), eventually becoming circular
- Discretize by replacing time derivative with difference in time; smooth curvature with one (of many) curvatures
- •Repeatedly add a little bit of κN ("forward Euler method")

 $\frac{d}{dt}\gamma(s,t) = -\kappa(s,t)N(s,t)$ $\frac{\gamma_i^{t+1} - \gamma_i^t}{\tau} = -\kappa_i^t N_i^t$ $\Rightarrow \gamma_i^{t+1} = \gamma_i^t - \tau \kappa_i^t N_i^t$

smooth

discrete

Elastic Flow

- Basic idea: rather than shrinking length, try to reduce bending (curvature)
- •Objective is integral of squared curvature; elastic flow is then gradient flow on this objective
- Minimizers are called *elastic curves*
- •More interesting w/ constraints (e.g., endpoint positions & a tangents)

http://brickisland.net/cs177fa12/?p=320



 $E(\gamma) := \int_0^L \kappa(s)^2 \, ds$ $\frac{d}{dt}\gamma = -\nabla_{\gamma}E(\gamma)$



Isometric Elastic Flow

- Different way to smooth out a curve is to directly "shrink" curvature
- Discrete case: "scale down" turning angles, then use the fundamental theorem of discrete plane curves to reconstruct
- Extremely stable numerically; exactly preserves edge lengths
- •Challenge: how do we make sure closed curves remain closed?



From Crane et al, "Robust Fairing via Conformal Curvature Flow"



Elastic Rods

- For space curve, can also try to minimize both *curvature* and *torsion*
- •Both in some sense measure "non-straightness" of curve
- Provides rich model of *elastic rods*
- •Lots of interesting applications (simulating hair, laying cable, ...)



From Bergou et al, "Discrete Elastic Rods"



Untangling Knots

- Is a given curve "knotted?"
- Minimize elastic energy *and* penalize self-collision
- *Might* go to smoothest curve in same isotopy class



 $\int_0^L \int_0^L \frac{1}{|\gamma(s) - \gamma(t)|^2} - \frac{1}{d(s,t)^2} \, ds \, dt$ Möbius energy



Credit: Henrik Schumacher





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